

# Theory of Computation

## Final Exam, Spring 2004

1. (15 pts) True or False? (No penalty for wrong answer.)

- 1)  If  $A$  is not recursive and  $B$  is a regular language, then  $\{x\#y \mid x \in A, y \in B, \# \text{ is a symbol not used in } A \text{ and } B\}$  is not recursive.
- 2)   $\{M \mid M \text{ is a TM and } L(M) \text{ contains at least } 3800000 \text{ elements}\}$  is recursive
- 3)   $\{M \mid M \text{ is a TM and } L(M) \text{ contains at most } 3800000 \text{ elements}\}$  is recursive.
- 4)  If  $A$  and  $B$  are context-free languages, so is  $A-B$  (i.e., the difference of  $A$  and  $B$ ).
- 5)  Every regular language is in  $P$  (polynomial time).
- 6)   $A \leq_m B$  and  $B$  is regular, then  $A$  is regular. (Here  $\leq_m$  denotes the many-one reduction.)
- 7)   $\{\langle M \rangle \mid M \text{ is a DFA and } L(M) \text{ is finite}\}$  is decidable.
- 8)  If 3CNF-SAT is in  $P$ , then  $P=NP$ .
- 9)  Given a PDA  $M$ , and a regular expression  $R$ , 'Is  $L(R)=L(M)$ ?' is decidable.
- 10)   $\{w w^R \mid w \text{ is binary string and } w^R \text{ is the reverse of } w\}$  is in  $P$ .
- 11)   $\{\langle M \rangle \mid M \text{ is a Turing machine that accepts the string } 1011\}$  is not recursive.
- 12)  A countably infinite union of recursive sets (i.e.,  $\cup_{i=1}^{\infty} L_i$ , each  $L_i$  is recursive) is always r.e.
- 13)  The set of all languages over the alphabet  $\{0,1\}$  is uncountable.
- 14)  There is an NP algorithm to determine whether or not two context-free grammars generate the same language.
- 15)   $\{a^k b^k c^k d^k e^k f^k \mid k \geq 0\} (\subseteq \{a, b, c, d, e, f\}^*)$  can be accepted by a 2-tape DTM in  $O(n)$  time.

2. (15 pts) Show the following problem to be NP-complete:

*Instance:* an undirected graph  $G = (V, E)$ .

*Question:* Is there a clique of size  $\lceil n/2 \rceil$  (i.e., half of  $n$ ) in  $G$ , where  $n = |V|$ , i.e., the number of nodes in  $V$ ?

(Hint: You may use any of the NP-complete problems discussed in class to do the reduction.)

3. (10 pts) Let  $G$  be the following grammar ( $\lambda$  is the empty string)

$$\begin{aligned} S &\rightarrow aB \mid bA \\ A &\rightarrow a \mid aS \mid bAA \\ B &\rightarrow \lambda \mid b \mid bS \mid aBB \end{aligned}$$

Construct an equivalent CFG in Chomsky Normal Form.

4. (20 pts) Answer whether a language class is closed under an operation by filling in *yes (Y)*, *no (N)*, *don't know (?)*

No penalty for wrong answer. No proofs are required.

| Operation/<br>language class | Union                 | Intersection                     | complement                       | Concatenation         | Kleene Star<br>(i.e., *) |
|------------------------------|-----------------------|----------------------------------|----------------------------------|-----------------------|--------------------------|
| <i>recursive</i>             | <input type="radio"/> | <input type="radio"/>            | <input type="radio"/>            | <input type="radio"/> | <input type="radio"/>    |
| <i>r.e.</i>                  | <input type="radio"/> | <input type="radio"/>            | <input checked="" type="radio"/> | <input type="radio"/> | <input type="radio"/>    |
| <i>Context-free</i>          | <input type="radio"/> | <input checked="" type="radio"/> | <input checked="" type="radio"/> | <input type="radio"/> | <input type="radio"/>    |
| <i>NP</i>                    | <input type="radio"/> | <input type="radio"/>            | <input type="radio"/>            | <input type="radio"/> | <input type="radio"/>    |

5. (20 pts)

- State the pumping lemma for context-free languages.
- Prove that  $\{a^p \mid p = n^2, n \in \mathbb{N}\}$  is not context-free.
- Use the pumping lemma as a tool to design an algorithm to decide whether a given context-free grammar  $G$  generates a finite set (i.e.,  $L(G)$  is finite).

6. (12 pts) Choosing from among **(D) decidable**, **(U) undecidable**, **(?) unknown**, categorize each of the following decision problems. No proofs are required. No penalty for wrong answer.

| Problem /<br>Language Class        | Regular               | Context Free                     | recursive                        | r.e.                             |
|------------------------------------|-----------------------|----------------------------------|----------------------------------|----------------------------------|
| $L = \Sigma^*$ ?                   | <input type="radio"/> | <input checked="" type="radio"/> | <input checked="" type="radio"/> | <input checked="" type="radio"/> |
| $L = \emptyset$ ?                  | <input type="radio"/> | <input type="radio"/>            | <input checked="" type="radio"/> | <input checked="" type="radio"/> |
| $x \in L$ , for<br>arbitrary $x$ ? | <input type="radio"/> | <input type="radio"/>            | <input type="radio"/>            | <input checked="" type="radio"/> |

7. (8 pts) Prove formally that if  $A \leq_m B$ , and  $B \leq_m C$ , the  $A \leq_m C$ . That is, the  $\leq_m$  relation (i.e., the many-one reduction) is transitive.