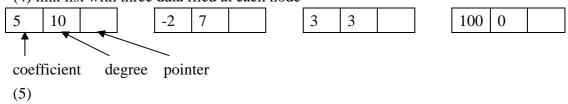
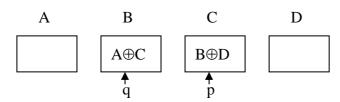
## **Data structures**

## **Fall 2003**

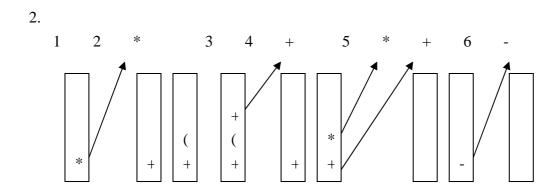
## **Midterm Solution**

- 1.
- (1) A collection of data along with specific operations that manipulate the data; It has nothing to do with a programming language. E.g. queue
- (2) Implementation of ADT. E.g. array
- (3) A BST is called balanced if the height of the tree is bounded by O(log n).
- (4) link list with three data filed at each node





traverse back:  $A = (content \ of \ B) \oplus p = (A \oplus C) \oplus C$ traverse forward:  $D = (content \ of \ C) \oplus q = (B \oplus D) \oplus B$ 



- 3. (1)  $\theta(1)$  (2)  $\theta(n)$  (3)  $\theta(\log n)$  (4)  $\theta(\log n)$  (5)  $\theta(n)$  (6)  $\theta(1)$  (7)  $\theta(n)$  (8)  $\theta(1)$  (9)  $\theta(1)$  (10)  $\theta(n)$
- 4.

height h. If an AVL tree with 50 nodes has height h it must be that  $m(h) \le 50 \le M(h)$ . Therefore, we are looking for the minimum and maximum h that satisfy the above inequalities.

As we have seen m(h) is given by the recurrence

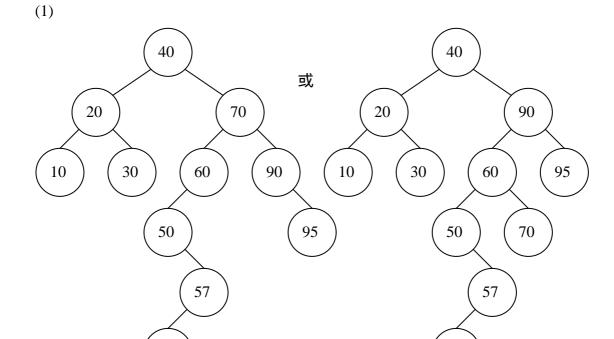
$$m(h) = \begin{cases} 1, & \text{if } h = 0 \\ 2, & \text{if } h = 1 \\ m(h-1) + m(h-2) + 1, & \text{if } h > 1 \end{cases}$$

On the other hand,  $M(h)=1+2+2^2+...+2^h=2^{h+1}-1$ . Therefore,

h	m(h)	M(h)		
0	1	1		
1	2	3		
2	4	7		
3	7	15		
4	12	31		
5	20	63		
6	33	127		
7	54	255		

So, the minimum and maximum h that satisfy  $m(h) \le 50 \le M(h)$  are 5 and 6.

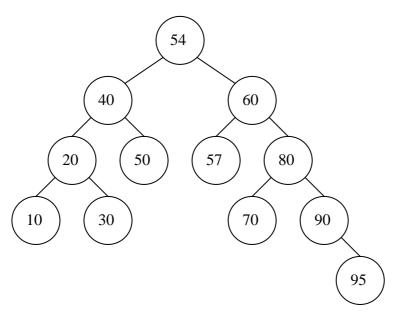
5.



54

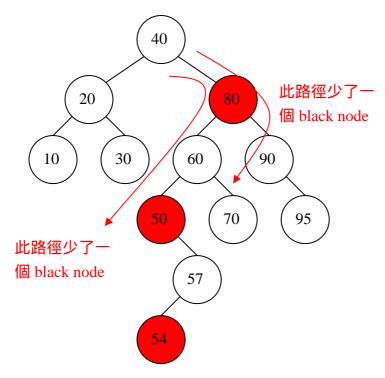
(2)

54



(3) preorder 40,20,10,30,80,60,50,57,54,70,90,95 inorder 10,20,30,40,50,54,57,60,70,80,90,95 postorder: 10,30,20,54,57,50,70,60,95,90,80,40

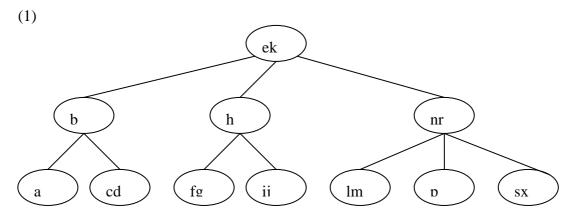
(4) 不行,因為根據 red-black tree 的定義,從 root 到任何 external node 的 path 要有相同多的 black nodes,但是在此題目中式不可能達到此要求,所以無法畫成 red-black tree



6.

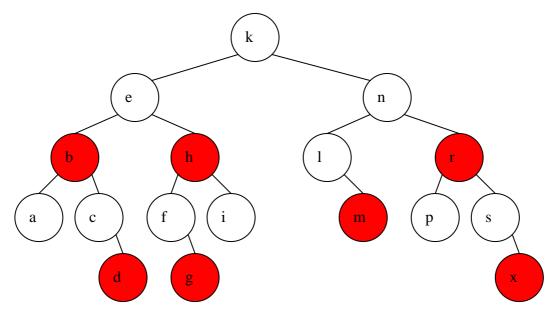
index	0	1	2	3	4	5	6	7	8	9
	20	0	42	99	75	15	25	36	47	89

7.



(2)

before insert j



8. 2B134BB